

# Supplementary Appendix

## *Risk Aversion in Share Auctions: Estimating Import Rents from TRQs in Switzerland.\**

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\*The main paper is available at <https://dx.doi.org/10.2139/ssrn.3397027>. Replicator files incl. data set are available at <https://github.com/SamuelHafner/RiskAversionInShareAuctions>.

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## A Proof of Proposition 1

The proof starts with an auxiliary lemma, which gives the interim utility for any opponent strategy profile  $\mu_{-i}$  and will be useful later on. To this end, I define

$$W_i(q|b_i, \mu_{-i}) = 1 - \int_{\mathcal{V}^{n-1}} \int_{\mathcal{B}^{n-1}} H_i^b(q) d\mu_1(b_1|v_1) \dots d\mu_{i-1}(b_{i-1}|v_{i-1}) d\mu_{i+1}(b_{i+1}|v_{i+1}) \dots d\mu_n(b_n|v_n) d\eta_{-i}(v_{-i}),$$

which returns the (decreasing) probability that the allocated quantity  $q_i^c$  for bidder  $i$  with valuation  $v_i$  strictly exceeds  $q \in [0, Q]$  when the submitted demand schedule is  $b_i$  and the opponent strategy profile is  $\mu_{-i}$ . Writing  $V_i(q) = \int_0^q v_i(q) dq$  and  $B_i(q) = \int_0^q \beta_{b_i}(q) dq$  for the respective gross valuation and gross payment accruing to bidder  $i$ , we have:

**Lemma A.1.** *Given an opponent strategy profile  $\mu_{-i} \in \mathcal{M}^{n-1}$ , the interim utility  $\Pi_i(b_i, v_i, \mu_{-i})$  for bidder  $i \in \{1, \dots, n\}$  of type  $v_i$  when submitting a bid schedule  $b_i \in \mathcal{B}$  is*

$$\Pi_i(b_i, v_i, \mu_{-i}) = \sum_{j=1}^k \int_{q^{j-1}}^{q^j} \phi'(V_i(q) - B_i(q)) [v_i(q) - p_i^j] W_i(q|b_i, \mu_{-i}) dq. \quad (\text{A.1})$$

*Proof of Lemma A.1.* Let

$$H_i^{(b_i, v_{-i})}(q) = \int_{\mathcal{B}^{n-1}} H_i^{(b_i, b_{-i})}(q) d\mu_1(b_1|v_1) \dots d\mu_{i-1}(b_{i-1}|v_{i-1}) d\mu_{i+1}(b_{i+1}|v_{i+1}) \dots d\mu_n(b_n|v_n) \quad (\text{A.2})$$

be the distribution of the quantity  $q_i^c$  that bidder  $i$  submitting  $b_i$  receives when his opponents play according to their strategies in  $\mu_{-i}$  and the opponent type profile is  $v_{-i}$ . Combining (1) and (A.2), the interim utility  $\Pi_i(b_i, v_i, \mu_{-i})$  of player  $i$  can be written as

$$\Pi_i(b_i, v_i, \mu_{-i}) = \int_{\mathcal{V}^{n-1}} \int_0^Q \phi(V_i(q) - B_i(q)) dH_i^{(b_i, v_{-i})}(q) d\eta_{-i}(v_{-i}). \quad (\text{A.3})$$

Because  $\phi(V_i(q) - B_i(q))$  is continuous and  $H_i^{(b_i, v_{-i})}(q)$  is monotone, the inner integral of the right-hand side in (A.3) can be integrated by parts (cf. Apostol, 1974, Theorem 7.6),

yielding

$$\begin{aligned} \Pi_i(b_i, v_i, \phi, \mu_{-i}) = \int_{\mathcal{V}^{n-1}} \left[ - \int_0^Q \phi'(V_i(q) - B_i(q)) [v_i(q) - \beta_{b_i}(q)] H_i^{(b_i, v_{-i})}(q) dq + \right. \\ \left. \phi(V_i(q) - B_i(q)) H_i^{(b_i, v_{-i})}(q) \Big|_0^Q \right] d\eta_{-i}(v_{-i}). \quad (\text{A.4}) \end{aligned}$$

We can rewrite (A.4) as

$$\begin{aligned} \Pi_i(b_i, v_i, \mu_{-i}) = - \int_{\mathcal{V}^{n-1}} \int_0^Q \phi'(V_i(q) - B_i(q)) [v_i(q) - \beta_{b_i}(q)] \times \\ H_i^{(b_i, v_{-i})}(q) dq d\eta_{-i}(v_{-i}) + \phi(V_i(Q) - B_i(Q)). \quad (\text{A.5}) \end{aligned}$$

Because  $f(v_{-i}, q) = \phi'(V_i(q) - B_i(q)) [v_i(q) - \beta_{b_i}(q)] H_i^{(b_i, v_{-i})}(q)$  is measurable and bounded on  $\mathcal{V}^{n-1} \times [0, Q]$ , the Fubini-Tonelli theorem (cf. [Rudin, 1970](#), Theorem 8.8) can be applied to get that (A.5) is

$$\begin{aligned} \Pi_i(b_i, v_i, \mu_{-i}) = - \int_0^Q \phi'(V_i(q) - B_i(q)) [v_i(q) - \beta_{b_i}(q)] \times \\ \int_{\mathcal{V}^{n-1}} H_i^{(b_i, v_{-i})}(q) d\eta_{-i}(v_{-i}) dq + \phi(V_i(Q) - B_i(Q)). \quad (\text{A.6}) \end{aligned}$$

By definition we have

$$W_i(q|b_i, \mu_{-i}) = 1 - \int_{\mathcal{V}^{n-1}} H_i^{(b_i, v_{-i})}(q) d\eta_{-i}(v_{-i}),$$

which allows to rewrite (A.6) as

$$\begin{aligned} \Pi_i(b_i, v_i, \mu_{-i}) = - \sum_{j=1}^{k+1} \int_{q_i^{j-1}}^{q_i^j} \phi'(V_i(q) - B_i(q)) [v_i(q) - p_i^j] (1 - W_i(q|b_i, \mu_{-i})) dq \\ + \phi(V_i(Q) - B_i(Q)), \end{aligned}$$

yielding (A.1), because  $W_i(q|b_i, \mu_{-i}) = 0$  for all  $q \in (q_i^k, Q]$  holds by assumption.  $\square$

Now, for finite natural  $h$ , let  $\mathcal{B}_{i,h}$  be a discrete action space for bidder  $i$  defined as

$$\mathcal{B}_{i,h} = \left\{ \{p_i^j, q_i^j\}_{j=1,\dots,k} \in [P_{i,h} \times Q_h]^k : p_i^j \geq p_i^{j+1}, q_i^j \leq q_i^{j+1}, q_i^{k+1} = Q, p_i^{k+1} = 0 \right\},$$

where

$$P_{i,h} = \left\{ 0, \frac{i}{nh^2}, \frac{1}{h} \left[ \bar{p} + \frac{i}{nh} \right], \frac{1}{h} \left[ 2\bar{p} + \frac{i}{nh} \right], \dots, \frac{1}{h} \left[ (h-1)\bar{p} + \frac{i}{nh} \right] \right\}$$

$$Q_h = \left\{ \frac{Q}{h}, 2\frac{Q}{h}, 3\frac{Q}{h}, \dots, Q \right\}.$$

For  $h$  large enough we have  $\mathcal{B}_{i,h} \subset \mathcal{B}$ . Moreover,  $\lim_{h \rightarrow \infty} \mathcal{B}_{i,h} = \mathcal{B}$ . The strategy of the proof is to first establish existence of an equilibrium  $\mu_h^*$  in the restricted auction with bidder-specific action spaces  $\mathcal{B}_{i,h}$  and, second, to use these equilibria to construct a sequence  $\mu_h^*$  of equilibria whose limit  $\mu^*$  is an equilibrium of the unrestricted auction with action space  $\mathcal{B}$  for all bidders. For this approach it is crucial that ties cannot occur at positive prices by the construction of the action spaces  $\mathcal{B}_{i,h}$  (cf. the proof to Lemma A.3 below).

Let  $\mathcal{M}_{i,h}$  be the space of distributional strategies on  $\mathcal{B}_{i,h} \times \mathcal{V}$  for player  $i$ . The next Lemma, which is a direct application of [Milgrom and Weber \(1985\)](#), establishes existence of an equilibrium  $\mu_h^*$  in the restricted auction for any  $h \in \mathbb{N}_+$ .

**Lemma A.2.** *There is an equilibrium  $\mu_h^* \in \times_{i \in \{1, \dots, n\}} \mathcal{M}_{i,h}$  in the restricted auction for any  $h$ .*

*Proof.* It follows from the Helly's selection theorem ([Rudin, 1964](#)) that  $\mathcal{V}$  is compact, and thus complete and separable. Because the action spaces  $\mathcal{B}_{i,h}$  are finite, they are compact and condition (a) in Proposition 1 of [Milgrom and Weber \(1985\)](#) is satisfied. Together with the type space assumption (A1) the assumptions of Theorem 1 in [Milgrom and Weber \(1985\)](#) are thus satisfied, and we have existence of an equilibrium that we denote by  $\mu_h^*$ .  $\square$

Having existence of an equilibrium  $\mu_h^* \in \times_{i \in \{1, \dots, n\}} \mathcal{M}_{i,h}$  allows me to show existence of an equilibrium with a distinct structure that we will need in the following. For  $q \in Q_h$ , let

$$\theta_i(q; v_i) = \max \{p \in P_{i,h} : p \leq v_i(q - Q/h)\},$$

and let  $\bar{\mathcal{M}}_{i,h} \subset \mathcal{M}_{i,h}$  be the set of strategies on

$$\{(b_i, v_i) \in \mathcal{B}_{i,h} \times \mathcal{V} : \beta_{b_i}(q) \leq \theta_i(q; v_i), \forall q \in Q_h\}.$$

That is, the support of the strategies in  $\bar{\mathcal{M}}_{i,h}$  consists of bids such that, in the limit  $h \rightarrow \infty$ , the corresponding step functions lie weakly below the marginal valuation function. For

further reference, let  $\bar{\mathcal{M}} = \lim_{h \rightarrow \infty} \bar{\mathcal{M}}_{i,h} \subseteq \mathcal{M}$  be the strategy space in this limit, which is independent of the bidder's identity, because  $\lim_{h \rightarrow \infty} \mathcal{B}_{i,h} = \mathcal{B}$  as observed above.

**Lemma A.3.** *An equilibrium  $\mu_h^*$  satisfying  $\mu_h^* \in \times_{i \in \{1, \dots, n\}} \bar{\mathcal{M}}_{i,h}$  exists for any  $h$ .*

*Proof.* Fix an equilibrium  $\mu_h^* \in \times_{i \in \{1, \dots, n\}} \mathcal{M}_{i,h}$ . I first argue that it is without loss of generality to assume that all bids  $b_i$  in the support of the equilibrium strategy of any bidder  $i$  with valuation  $v_i \in V$  satisfy

$$q_i^j \leq q_i^{j+1} - Q/h \text{ whenever } p_i^j > 0$$

for all  $j \in \{1, \dots, k-1\}$ . To see this, consider a bid  $b_i$  for which this is not the case, that is, suppose we have  $b_i$  with  $q_i^j = q_i^{j+1}$  for at least one  $j \in \{1, \dots, k-1\}$ . As the price-quantity-pairs  $(p_i^{j+1}, q_i^{j+1})$  will be ignored by the auctioneer and there is zero probability to win a quantity  $q$  for which the price bid is zero, such a bid  $b_i$  yields the same payoff as a bid  $b'_i$  that is equal to  $b_i$  except that the price-quantity-pairs  $(p_i^{j+1}, q_i^{j+1})$  are deleted and there is an equal number of price-quantity pairs  $(0, Q)$  appended at the end. For example, if there is one such pair  $j$ , then we have

$$b'_i = \{(p_i^1, q_i^1), \dots, (p_i^j, q_i^j), (p_i^{j+2}, q_i^{j+2}), \dots, (p_i^k, q_i^k), (0, Q)\}.$$

Any bid  $b'_i$  that is thus altered does not change the utility of the other players, and hence for any equilibrium strategy  $\mu_{i,h}^*$  having  $b_i$  in its support there is an alternative strategy  $\mu'_{i,h}$  constructed from  $\mu_{i,h}^*$  with all the mass on  $b_i$  appropriately shifted to  $b'_i$ , so that  $\mu_h^* = (\mu'_{i,h}, \mu_{-i}^*)$  is also an equilibrium.

So, suppose  $b_i = \{(p_i^1, q_i^1), \dots, (p_i^k, q_i^k)\} \in \mathcal{B}_{i,h}$  with  $q_i^j \leq q_i^{j+1} - Q/h$  for  $j \in \{1, \dots, \ell_i - 1\}$  where  $\ell_i \leq k$ , and, if  $\ell_i < k$ , with  $(p_i^j, q_i^j) = (0, Q)$  for  $j \in \{\ell_i + 1, \dots, k\}$  is in the support of the equilibrium strategy of bidder  $i$  with valuation  $v_i$ . I now show that  $\beta_{b_i}(q) \leq \theta_i(q; v_i)$  holds for all  $q \in Q_h$ . By optimality, we get from interim utility (A.1) that is holds, for any  $j \in \{1, \dots, \ell_i\}$  (taking  $q_i^{k+1} = Q$  and  $p_i^{k+1} = 0$ ),

$$\begin{aligned} & \int_{q_i^{j-1}}^{q_i^j} \phi'(V_i(q) - B_i(q)) [v_i(q) - p_i^j] W_i^j(q|b_i, \mu_{-i}) dq \\ & + \int_{q_i^j}^{q_i^{j+1}} \phi'(V_i(q) - B_i(q)) [v_i(q) - p_i^{j+1}] W_i^{j+1}(q|b_i, \mu_{-i}) dq \geq \\ & \int_{q_i^{j-1}}^{q_i^j - Q/h} \phi'(V_i(q) - B_i(q)) [v_i(q) - p_i^j] W_i^j(q|b'_i, \mu_{-i}) dq \\ & + \int_{q_i^j - Q/h}^{q_i^{j+1}} \phi'(V_i(q) - B_i(q)) [v_i(q) - p_i^{j+1}] W_i^{j+1}(q|b'_i, \mu_{-i}) dq, \quad (\text{A.7}) \end{aligned}$$

where  $b'_i$  is equal to  $b_i$  except for the  $j$ -th quantity point,  $q_i^j$ , which is replaced by  $q_i^j - Q/h$ .

Because ties at positive prices cannot happen and no quantities for which the price bid is zero are ever allocated, the probability to win a certain quantity only depends on the price bid for that quantity. That is, we have  $W_i^j(q|b_i, \mu_{-i}) = W_i^j(q|b'_i, \mu_{-i})$  on  $[q_i^{j-1}, q_i^j - Q/h]$  and  $W_i^{j+1}(q|b_i, \mu_{-i}) = W_i^{j+1}(q|b'_i, \mu_{-i})$  on  $[q_i^j, q_i^{j+1}]$ . Consequently, it follows from (A.7) that

$$\begin{aligned} \int_{q_i^j - Q/h}^{q_i^j} \phi' (V_i(q) - B_i(q)) [v_i(q) - p_i^j] W_i^j(q|b_i, \mu_{-i}) dq \\ \geq \int_{q_i^j - Q/h}^{q_i^j} \phi' (V_i(q) - B_i(q)) [v_i(q) - p_i^{j+1}] W_i^{j+1}(q|b'_i, \mu_{-i}) dq. \end{aligned} \quad (\text{A.8})$$

First, consider step  $j = \ell_i$ . Because no quantities for which a prize of zero is bid are ever won, condition (A.8) becomes

$$\int_{q_i^{\ell_i} - Q/h}^{q_i^{\ell_i}} \phi' (V_i(q) - B_i(q)) [v_i(q) - p_i^{\ell_i}] W_i^{\ell_i}(q|b_i, \mu_{-i}) dq \geq 0,$$

which together with the assumptions that  $\phi'(\cdot) > 0$  and that the marginal value  $v_i$  is decreasing gives us that

$$p_i^{\ell_i} \leq v_i(q_i^{\ell_i} - Q/h) \quad (\text{A.9})$$

must hold.

But with inequality (A.9) at hand, I can now argue that  $p_i^j \leq v_i(q_i^j - Q/h)$  must hold for every  $j \in \{1, \dots, \ell_i\}$ : Suppose, to the contrary, that it does not hold for some  $j < \ell_i$ ; i.e., we have  $p_i^j > v_i(q_i^j - Q/h)$ . Because  $v_i$  is decreasing, this implies that the left side of (A.8) is strictly negative, so that it must hold that  $p_i^{j+1} > v_i(q_i^j)$ , which is necessary for right side of (A.8) to be strictly negative, too. But because  $q_i^j \leq q_i^{j+1} - Q/h$  holds and  $v_i$  is decreasing,  $p_i^{j+1} > v_i(q_i^j)$  in turn implies that  $p_i^{j+1} > v_i(q_i^{j+1} - Q/h)$  holds, as well. Repeating this for every  $j' > j$ , ultimately yields  $p_i^{\ell_i} > v_i(q_i^{\ell_i} - Q/h)$ , which contradicts (A.9). As we can apply this argument to any  $j < \ell_i$ , we get that

$$p_i^j \leq v_i(q_i^j - Q/h), \quad \forall j \in \{1, \dots, k\}$$

must hold for any  $v_i$ , thus giving us the claim.  $\square$

Next, consider a sequence of auctions with restricted action space  $\times_{i \in \{1, \dots, n\}} \mathcal{B}_{i,h}$  having equilibria  $\mu_h^* \in \times_{i \in \{1, \dots, n\}} \bar{\mathcal{M}}_{i,h}$  for  $h \rightarrow \infty$ . Because the space  $\mathcal{M}$  of probability measures on  $\mathcal{B} \times \mathcal{V}$  is compact in the weak\*-topology (Milgrom and Weber, 1985), the sequence  $\mu_h^* \in \times_{i \in \{1, \dots, n\}} \bar{\mathcal{M}}_{i,h} \subset \mathcal{M}^n$  has a converging subsequence. Pick such a subsequence and suppose it

converges to some  $\mu^*$ . By Remark 3.1 in [Reny \(1999\)](#)  $\mu^*$  is an equilibrium in the unrestricted auction if the unrestricted auction with strategy space  $\bar{\mathcal{M}}^n$  is better reply secure and if for every  $\epsilon > 0$  there is  $H > 0$  such that for all  $h > H$  the profile  $\mu_h^*$  in the respective subsequence is an  $\epsilon$ -equilibrium of the unrestricted auction. In the following I write  $U_i(\mu) = \int \Pi_i(b_i, v_i, \mu_{-i}) d\mu_i$  for player  $i$ 's payoff function, where  $\Pi_i$  is player  $i$ 's interim utility as defined in Lemma [A.1](#).

**Definition 1** (Better-Reply Security, cf. [Reny 1999](#)). *Game  $G = (\bar{\mathcal{M}}, U_i)_{i \in \{1, \dots, n\}}$  where each player  $i = 1, \dots, n$  has a strategy space  $\bar{\mathcal{M}}$  and a payoff function  $U_i(\mu)$ ,  $\mu \in \bar{\mathcal{M}}^n$ , is better-reply secure if whenever  $(\mu^*, u^*)$  is in the closure of the graph of its vector payoff function  $U$ ,  $\{(\mu, u) : u = U(\mu)\}$ , and  $\mu^*$  is not a Nash equilibrium, then some player  $i$  can secure a payoff strictly above  $u_i^*$  at  $\mu^*$ : There exists some  $\tilde{\mu}_i$  such that  $U_i(\tilde{\mu}_i, \mu_{-i}) > u_i^*$  for all  $\mu_{-i}$  in some open neighborhood of  $\mu_{-i}^*$ .*

**Definition 2** ( $\epsilon$ -equilibrium). *A strategy profile  $\mu \in \bar{\mathcal{M}}^n$  is an  $\epsilon$ -equilibrium of game  $G = (\bar{\mathcal{M}}, U_i)_{i \in N}$  if for all players  $i \in N$  it holds  $U_i(\hat{\mu}_i, \mu_{-i}) - U_i(\mu) \leq \epsilon$  for every  $\hat{\mu}_i \in \bar{\mathcal{M}}$ .*

*Step I: Better-Reply Security.* I start by showing better reply security, for which I adapt the argument given in [Reny \(1999\)](#) for the multi-unit auction case. For opponent profile  $\mu_{-i} \in \bar{\mathcal{M}}^{n-1}$ , let

$$B^\epsilon(\mu_{-i}) = \left\{ \mu_i \in \bar{\mathcal{M}} : \left| U_i(\mu_i, \mu_{-i}) - \sup_{\tilde{\mu}_i \in \bar{\mathcal{M}}} U_i(\tilde{\mu}_i, \mu_{-i}) \right| \leq \epsilon \right\}$$

be the set of strategies  $\mu_i$  that yield utility within  $\epsilon > 0$  of the supremum. The following observation is needed below.

**Lemma A.4.** *Fix  $\tilde{\mu}_{-i} \in \bar{\mathcal{M}}^{n-1}$ . Then, for every  $\epsilon > 0$  sufficiently small and for any  $\mu_i \in B^\epsilon(\tilde{\mu}_{-i})$ ,  $U_i(\mu_i, \mu_{-i})$  is continuous in  $\mu_{-i}$  at  $(\mu_i, \tilde{\mu}_{-i})$ .*

*Proof.* By contradiction. Take  $\mu_i \in B^\epsilon(\tilde{\mu}_{-i})$  and suppose  $U_i(\mu_i, \mu_{-i})$  is not continuous in  $\mu_{-i}$  at  $(\mu_i, \tilde{\mu}_{-i})$ . If  $U_i(\mu_i, \mu_{-i})$  is not continuous in  $\mu_{-i}$  at  $(\mu_i, \tilde{\mu}_{-i})$ , then there must be a bidder  $j \in \{1, \dots, n\} \setminus i$  and a clearing price  $p^c$  such that bidder  $i$  and bidder  $j$  tie at  $p^c$  with positive probability. That is, there are  $X, Y \subset \mathcal{V}$  with  $\eta_i(X), \eta_j(Y) > 0$  such that both bidders have price points  $p_i^{m_i} = p^c$  and  $p_j^{m_j} = p^c$  (where the steps  $m_i$  and  $m_j$  might be distinct for the two bidders) in the support of their strategies  $\mu_i(\cdot | v_i)$  and  $\mu_j(\cdot | v_j)$  whenever they are of a type  $v_i \in X$  and  $v_j \in Y$ , respectively.

The discontinuity together with the fact that  $\sum_{i \in N} q_i^c = Q$  implies that at least one tying bidder is rationed with positive probability. Without loss suppose this to be bidder  $i$ . From Assumption [\(A4\)](#) we then obtain that the expected allocated quantity lies in

$[q_i^{m_i-1}, \min\{q_i^j, Q - \lim_{p \searrow p^c} \sum_i \beta_{b_i}^{-1}(p)\})$ . Moreover, (A4) implies that bidder  $i$  could secure a quantity of at least  $\min\{q_i^j, Q - \lim_{p \searrow p^c} \sum_i \beta_{b_i}^{-1}(p)\}$  by marginally raising the price point  $p_i^{m_i}$ . Now, recall that it follows from (A2) that there are  $X', X'' \subset X_1$  with  $\eta_i(X'), \eta_i(X'') > 0$  where  $\forall f \in X'$  and  $\forall g \in X''$  it holds that  $f(q) > g(q), \forall q \in [0, 1]$ . Because  $\phi$  is strictly increasing, this gives us that there is indeed a set of bidder  $i$  types with strictly positive measure that strictly prefer to avoid the tie by marginally raising  $p_i^{m_i}$ . As the increase in utility is strict, we have, for any  $\epsilon > 0$  sufficiently small, a contradiction to the assumption that  $\mu_i \in B^\epsilon(\tilde{\mu}_{-i})$ .  $\square$

Next, consider some  $(\mu^*, u^*)$  which is in the closure of the graph of the payoff function (i.e., there is a sequence  $\mu_m \rightarrow \mu^*$  such that  $u^* = \lim_{m \rightarrow \infty} U(\mu_m)$ ) without  $\mu^*$  being an equilibrium. To show better-reply security, we need to establish that there is some bidder  $i$  that can secure a payoff strictly above  $u_i^*$  by deviating from  $\mu_i^*$  even if the other bidders also slightly deviate. Two cases need to be considered: (i)  $U(\cdot)$  is continuous at  $\mu^*$  and (ii)  $U(\cdot)$  is not continuous at  $\mu^*$ .

- (i) Consider first the case of  $U(\cdot)$  being continuous at  $\mu^*$ . Then there is a bidder  $i$ , an  $\epsilon > 0$  small enough, and some  $\mu_i \in B^\epsilon(\mu_{-i}^*)$  such that  $U_i(\mu_i, \mu_{-i}^*) > U(\mu^*) = u^*$ . As  $U_i(\mu_i, \cdot)$  is continuous at  $\mu_{-i}^*$  by Lemma A.4, we have better reply security.
- (ii) Second, consider the case of  $U(\cdot)$  being discontinuous at  $\mu^*$ : There must be at least two bidders  $i$  and  $j$ , a clearing price  $p^c$ , and a sequence  $\mu_m \rightarrow \mu^*$  such that both bidders have a positive measure of types that have price points  $p_i^{t_i} = p^c$  and  $p_j^{t_j} = p^c$  (where the steps  $t_i$  and  $t_j$  might be distinct for the two bidders) in the support of their strategies in the limit  $\mu^*$  but do not tie at  $p^c$  for any  $\mu_m$  along the sequence whenever  $m$  is sufficiently high. Without loss, suppose  $p_i^{t_i} < p_j^{t_j}$  for any  $\mu_m$  with sufficiently high  $m$ .

To continue, observe that there is some sufficiently small  $\epsilon > 0$  and some  $\hat{\mu}_i \in B^\epsilon(\mu_{-i}^*)$  such that  $U_i(\hat{\mu}_i, \mu_{-i}^*) > U_i(\mu^*)$ , which, by the same argument as in the proof to Lemma A.4, involves a positive mass of types marginally increasing the price point which was equal to the clearing price  $p^c$ . But then,  $\hat{\mu}_i$  yields for these types a strictly higher utility than  $\mu_{i,m}$  does against  $\mu_{-i,m}$  for any sufficiently high  $m$  (because by marginally raising the respective price point discontinuously raises the winning probability for the respective quantities for all sufficiently high  $m$ ); i.e. there is  $\delta > 0$  and  $M > 0$  such that

$$U_i(\hat{\mu}_i, \mu_{-i,m}) - U_i(\mu_m) > \delta, \forall m > M.$$



Because  $U_i(\hat{\mu}_i, \mu_{-i})$  is continuous in  $\mu_{-i}$  at  $\mu_{-i}^*$  by Lemma A.4, it follows that

$$U_i(\hat{\mu}_i, \mu_{-i}^*) > \lim_{m \rightarrow \infty} U_i(\mu_m) = u_i^*,$$

giving us better-reply security in this case, too.

*Step II:  $\epsilon$ -Equilibria.* I follow [Reny \(2011\)](#) and show that for every  $\epsilon > 0$  there is  $h$  high enough such that for every feasible action  $b_i$  there is a feasible action  $b_{i,h}$  such that the ex-post loss from choosing  $b_{i,h}$  rather than  $b_i$  is smaller than  $\epsilon$ , and that this holds uniformly in the strategies  $\mu_{-i}$  of the other players.

Fix some finite natural  $h$ , some bidder  $i$  with type  $(v_i, \phi)$  and any  $b_i$  for which it holds that  $\beta_{b_i}(q) \leq \theta_i(q; v_i)$ ,  $\forall q \in Q_h$ . If  $b_i \in \mathcal{B}_{i,h}$  then we are done. So consider  $b_i \notin \mathcal{B}_{i,h}$ . Let

$$b_i = \{(p_i^1, q_i^1), \dots, (p_i^k, q_i^k)\},$$

and define

$$b_{i,h} = \{(p_{i,h}^1, q_{i,h}^1), \dots, (p_{i,h}^k, q_{i,h}^k)\},$$

with

$$\begin{aligned} p_{i,h}^j &= \min \{p \in P_{i,h} : p \geq p_i^j\} \\ q_{i,h}^j &= \max \{q \in Q_h : q \leq q_i^j \text{ and } p_{i,h}^j \leq v_i(q - Q/h)\}, \end{aligned}$$

for all  $j \in \{1, \dots, k\}$ . Above definitions guarantee that  $\beta_{b_{i,h}}(q) \leq \theta_i(q, v_i)$  holds for all  $q \in Q_h$ , and hence that  $b_{i,h}$  is a feasible action, as well as that  $q_{i,h}^j \rightarrow q_i^j$ . The ex-post loss sources from switching from  $b_i$  to  $b_{i,h}$  are threefold:

1. There might be shares  $q$  at which it holds that  $\beta_{b_{i,h}}(q) > \beta_{b_i}(q)$  and that are won under  $b_i$ , that are also won under  $b_{i,h}$  yet at a higher price. The loss from such quantities is bounded above by

$$\phi(V_i(q_i^k) - B_i(q_i^k)) - \phi\left(V_i(q_i^k) - B_i(q_i^k) - \sum_{j=1}^k (q_i^j - q_i^{j-1}) \frac{\bar{p}}{h}\right). \quad (\text{A.10})$$

2. There might be shares  $q$  at which it holds that  $\beta_{b_{i,h}}(q) > v_i(q) \geq \beta_{b_i}(q)$  and that are not won under  $b_i$ , but that are won under  $b_{i,h}$ . The loss from such quantities is also bounded above by

$$\phi(V_i(q_i^k) - B_i(q_i^k)) - \phi\left(V_i(q_i^k) - B_i(q_i^k) - \sum_{j=1}^k (q_i^j - q_i^{j-1}) \frac{\bar{p}}{h}\right). \quad (\text{A.11})$$

3. There might be shares  $q$  at which it holds that  $\beta_{b_{i,h}}(q) < \beta_{b_i}(q)$  and that are won under  $b_i$ , but that are not won under  $b_i^h$ . The loss from such quantities is bounded above by

$$\sum_{j=1}^k \int_{q_{i,h}^j}^{q_i^j} \phi'(V_i(q) - B_i(q)) v_i(q) dq. \quad (\text{A.12})$$

All three bounds (A.10)–(A.12) vanish as  $h \rightarrow \infty$  independently, and hence uniformly, in the strategies  $\mu_{-i}$  of the other players, because  $\phi'(\cdot)$  is bounded on  $\mathbb{R}$ . Consequently, we have that  $\mu_h^*$  is a sequence of  $\epsilon$ -equilibria of the unrestricted auction where  $\epsilon \rightarrow 0$  when  $h \rightarrow \infty$ , as required.

We can conclude that an equilibrium exists. Moreover, absence of ties follows from the same argument as in the proof of Lemma A.4: If a tie were to happen with positive probability then there would always be a non-negligible set of types for at least one of the bidders that strictly prefer to avoid the tie.

## B Inequality Violations By Bidder and Auction Groups

This appendix provides more details on the shares of inequality violations in the respective bidder groups and the respective auction groups that I use for estimation.

### B.1 Bidder Groups

The tables in Table 1 report the full set of the estimated values of  $\Theta_g(\rho)$  for the three bidder groups  $g = 1, 2, 3$ . These tables complement the tables presented in Figure 3 in Section 5.2 of the main text.

**Table 1:** *The three tables describe the estimates of  $\Theta_g(\rho)$  for the three different bidder groups that I use for estimation.*

Group $g = 1$			Group $g = 2$			Group $g = 3$		
$\rho$	$\Theta_g$	se	$\rho$	$\Theta_g$	se	$\rho$	$\Theta_g$	se
0.0	0.527	0.01473	0.0	0.5014	0.01414	0.0	0.4885	0.02084
$4.54e - 5$	0.4836	0.01349	$4.54e - 5$	0.343	0.01294	$4.54e - 5$	0.2982	0.01787
$7.485e - 5$	0.4606	0.01303	$7.485e - 5$	0.2987	0.01222	$7.485e - 5$	0.2742	0.01714
0.0001234	0.4295	0.01231	0.0001234	0.2547	0.01139	0.0001234	0.2585	0.01637
0.0002035	0.3915	0.01132	0.0002035	0.2152	0.01067	0.0002035	0.2448	0.01643
0.0003355	0.349	0.01043	0.0003355	0.1853	0.01115	0.0003355	0.2443	0.01728
0.0005531	0.3041	0.009707	0.0005531	0.166	0.01186	0.0005531	0.2563	0.01719
0.0009119	0.2615	0.009021	0.0009119	0.1589	0.01255	0.0009119	0.2798	0.01631
0.001503	0.2244	0.008897	0.001503	0.1671	0.01292	0.001503	0.3009	0.0148
0.002479	0.1954	0.009447	0.002479	0.1922	0.01315	0.002479	0.3298	0.01549
0.004087	0.178	0.009966	0.004087	0.2348	0.01377	0.004087	0.3767	0.01826
0.006738	0.1748	0.01061	0.006738	0.2916	0.01394	0.006738	0.4338	0.01759
0.01111	0.1868	0.01129	0.01111	0.3627	0.01522	0.01111	0.4835	0.01785
0.01832	0.2136	0.01218	0.01832	0.443	0.01578	0.01832	0.5391	0.01973
0.0302	0.2563	0.01236	0.0302	0.5236	0.01802	0.0302	0.5795	0.01789
0.04979	0.3093	0.0123	0.04979	0.5971	0.01907	0.04979	0.606	0.01612
0.08208	0.3724	0.01351	0.08208	0.6512	0.01687	0.08208	0.6247	0.01551
0.1353	0.4445	0.0152	0.1353	0.6836	0.01383	0.1353	0.6396	0.01473
0.2231	0.5181	0.01575	0.2231	0.7009	0.01175	0.2231	0.6507	0.01375
0.3679	0.5806	0.01495	0.3679	0.7089	0.01016	0.3679	0.6605	0.01276
0.6065	0.6296	0.0144	0.6065	0.7131	0.009502	0.6065	0.6685	0.01188
1.0	0.6649	0.01312	1.0	0.7163	0.009106	1.0	0.6749	0.0116

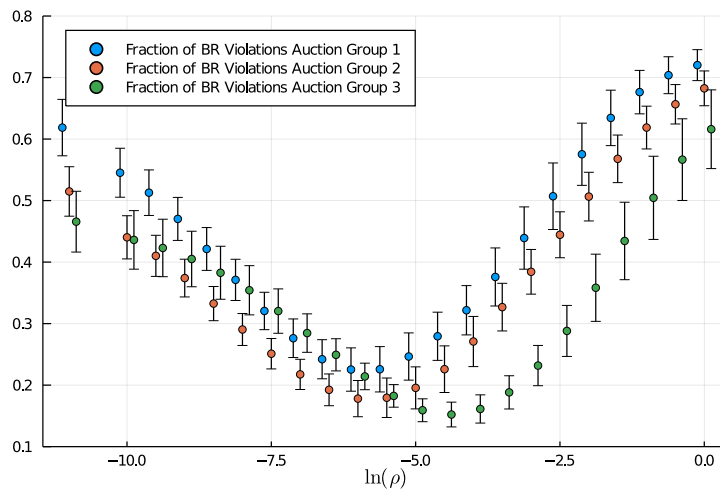
## B.2 Auction Groups

Figure 1 and Table 2 compare the shares of inequality violations between the three different auction groups that I use for estimation (cf. Appendix B of the main text). For simplicity I assume that all bidders have the same risk-aversion parameter  $\rho$ . Specifically, let  $\mathcal{T}_a \subset \{1, \dots, T\}$  be the set of auctions in auction group  $a = 1, 2, 3$ . For every of these auction groups, I compute

$$\Theta_a(\rho) \equiv \frac{1}{|\mathcal{T}_a|} \sum_{t \in \mathcal{T}_a} \left[ \frac{1}{\sum_{i \in N_t} \ell_{i,t}} \left[ \sum_{i \in N_t} \sum_{j=1}^{\ell_{i,t}} [\mathbf{1}\{(F_i^j(b_i, \hat{v}_i, \rho) > 0) \text{ or } (F_i^j(b_i, \hat{v}_i, \rho) < 0)\}] \right] \right].$$

The estimate  $\Theta_a(\rho)$  is U-shaped for all three groups. Ideally, the minimizers of this function are the same for all auction groups  $a = 1, 2, 3$ , which would indicate that risk preferences remain stable across auctions. Yet, the values of  $\rho$  that minimize the estimated  $\Theta_a(\rho)$  differ between groups 1 – 2, for which the minimum is at  $\rho = 0.0025$ , and group 3, for which the minimum is at  $\rho = 0.0111$ .

Nevertheless, these estimates do not allow to conclude that the assumption of stable risk preferences does not hold. To see this, consider the value  $\rho = 0.0041$ , corresponding to  $\ln(\rho) = -4.5$ . (The value  $\rho = 0.0041$  minimizes the fraction of inequality violations when assuming the same risk preference across all bidders and auctions; cf. the next appendix.) For each group  $a = 1, 2, 3$ , the 95% confidence interval around the estimated value of  $\Theta_a$  at  $\rho = 0.0041$  overlaps with the confidence interval around the respective minimum of  $\Theta_a$ . In other words, for neither group  $a = 1, 2, 3$  it is the case that the share of inequality violations at the minimum of  $\Theta_a$  significantly differs from that at  $\rho = 0.0041$ .



**Figure 1:** The figure shows the fraction of inequality violations,  $\Theta(\rho)$ , for the different auction groups used for estimation.

**Table 2:** *The three tables describe the estimates of  $\Theta_a(\rho)$  for the three different auction groups that I use for estimation.*

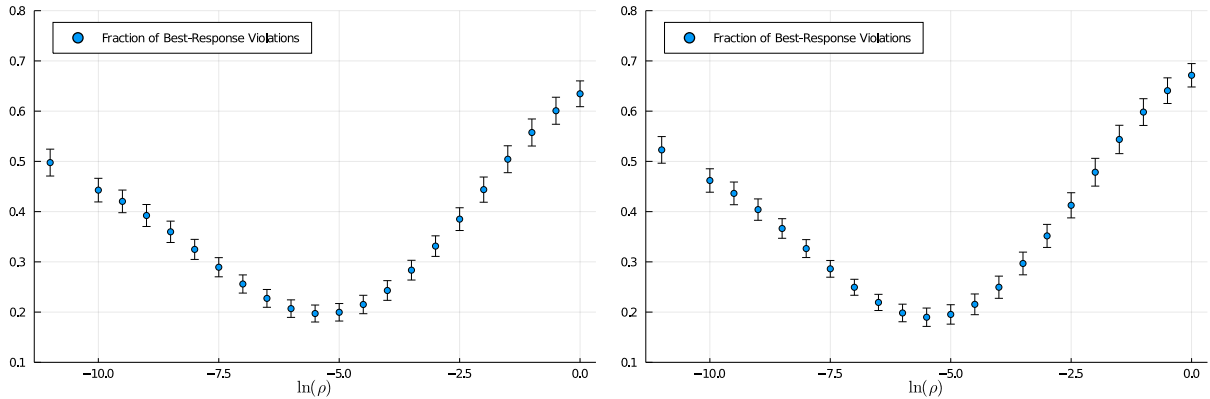
Group $a = 1$			Group $a = 2$			Group $a = 3$		
$\rho$	$\Theta_a$	se	$\rho$	$\Theta_a$	se	$\rho$	$\Theta_a$	se
0.0	0.6187	0.0234	0.0	0.5148	0.02052	0.0	0.4657	0.02521
$4.54e - 5$	0.5453	0.0203	$4.54e - 5$	0.4403	0.01791	$4.54e - 5$	0.4361	0.02428
$7.485e - 5$	0.5128	0.01889	$7.485e - 5$	0.4102	0.01705	$7.485e - 5$	0.423	0.02384
0.0001234	0.4702	0.01784	0.0001234	0.374	0.01565	0.0001234	0.405	0.02303
0.0002035	0.4213	0.01776	0.0002035	0.3325	0.01418	0.0002035	0.3827	0.02198
0.0003355	0.3711	0.01709	0.0003355	0.2904	0.01324	0.0003355	0.3542	0.0204
0.0005531	0.3206	0.01551	0.0005531	0.2511	0.0126	0.0005531	0.3204	0.01841
0.0009119	0.2762	0.01601	0.0009119	0.2174	0.01252	0.0009119	0.2846	0.01596
0.001503	0.2421	0.01622	0.001503	0.1924	0.01318	0.001503	0.2493	0.01334
0.002479	0.2253	0.01797	0.002479	0.1781	0.01502	0.002479	0.2141	0.01101
0.004087	0.2258	0.0188	0.004087	0.1795	0.01637	0.004087	0.1825	0.00936
0.006738	0.2465	0.01959	0.006738	0.1955	0.01743	0.006738	0.1592	0.009487
0.01111	0.2794	0.01999	0.01111	0.2259	0.01934	0.01111	0.1522	0.01027
0.01832	0.3217	0.02044	0.01832	0.2709	0.0208	0.01832	0.1613	0.0117
0.0302	0.3758	0.02408	0.0302	0.3268	0.01976	0.0302	0.1882	0.01378
0.04979	0.4391	0.02584	0.04979	0.3842	0.01853	0.04979	0.2318	0.01669
0.08208	0.507	0.0276	0.08208	0.4444	0.01905	0.08208	0.2881	0.02115
0.1353	0.5753	0.02579	0.1353	0.5065	0.02018	0.1353	0.3583	0.02786
0.2231	0.6344	0.02299	0.2231	0.5679	0.01973	0.2231	0.4344	0.03212
0.3679	0.6763	0.01802	0.3679	0.6187	0.01774	0.3679	0.5045	0.03449
0.6065	0.7038	0.01532	0.6065	0.6566	0.01635	0.6065	0.5666	0.03387
1.0	0.7202	0.01284	1.0	0.6825	0.01444	1.0	0.616	0.03264

## C An alternative functional form for the cdf of $D_{g,t}(p)$

In this appendix, I analyze an alternative assumption about the distribution family for residual demand. Specifically, I assume that residual demand  $D_{g,t}(p)$  follows a log-normal distribution on  $(0, \infty)$  (cf. Assumption (A6) in the main text). To determine how such an assumption compares to that of the gamma distribution used in the main analysis, I compute the fraction of best-response violations under the assumption that all bidders have the same risk-aversion parameter  $\rho$ . Specifically, I compute

$$\Theta(\rho) \equiv \frac{1}{T} \sum_{t=1}^T \left[ \frac{1}{\sum_{i \in N_t} \ell_{i,t}} \left[ \sum_{i \in N_t} \sum_{j=1}^{\ell_{i,t}} [\mathbf{1}\{(F_i^j(b_i, \hat{v}_i, \rho) > 0) \text{ or } (F_i^j(b_i, \hat{v}_i, \rho) < 0)\}] \right] \right].$$

Figure 2 compares the estimate  $\Theta(\rho)$  for the log-normal (left panel) and the gamma distribution (right panel). Under either assumptions,  $\Theta$  is U-shaped and the lowest value of  $\Theta(\rho)$  is at  $\rho = 0.0041$ . The numbers in Table 3 show that the fraction of best response violations are indeed comparable under the two assumptions. I conclude that the posited risk preference has a much higher impact on model fit than the specific assumption about the distribution of residual demand.



**Figure 2:** The left graph shows  $\Theta(\rho)$  when  $W_i^*$  and  $w_i^*$  are estimated under the assumption of a log normal distribution for  $D_{g,t}(p)$ . The right graph shows  $\Theta(\rho)$  under the assumption of a gamma distribution.

**Table 3:** The tables show the values of  $\Theta(\rho)$  when  $W_i^*$  and  $w_i^*$  are estimated under the assumption of a log normal distribution for  $D_{g,t}(\rho)$  and when assuming a gamma distribution (as in the main analysis).

Log Normal			Gamma		
$\rho$	$\Theta$	se	$\rho$	$\Theta$	se
0.0	0.4976	0.01365	0.0	0.5229	0.01353
$4.54e - 5$	0.4428	0.012	$4.54e - 5$	0.462	0.0119
$7.485e - 5$	0.4204	0.0115	$7.485e - 5$	0.4363	0.01153
0.0001234	0.3923	0.01116	0.0001234	0.404	0.01085
0.0002035	0.3598	0.01084	0.0002035	0.3665	0.009948
0.0003355	0.3248	0.01019	0.0003355	0.3264	0.009107
0.0005531	0.2893	0.009733	0.0005531	0.2861	0.008498
0.0009119	0.2559	0.009227	0.0009119	0.2494	0.008114
0.001503	0.2273	0.009051	0.001503	0.2193	0.008206
0.002479	0.2069	0.008902	0.002479	0.1983	0.008927
0.004087	0.1972	0.008625	0.004087	0.1898	0.009317
0.006738	0.1996	0.00889	0.006738	0.1954	0.009861
0.01111	0.2152	0.009352	0.01111	0.2155	0.01052
0.01832	0.243	0.01	0.01832	0.2495	0.01127
0.0302	0.2835	0.01003	0.0302	0.2968	0.01152
0.04979	0.3313	0.01044	0.04979	0.3517	0.01171
0.08208	0.3851	0.01156	0.08208	0.4126	0.01278
0.1353	0.4438	0.0128	0.1353	0.4784	0.01411
0.2231	0.5043	0.01367	0.2231	0.5437	0.01441
0.3679	0.5575	0.01373	0.3679	0.5982	0.01358
0.6065	0.6008	0.01373	0.6065	0.6408	0.01302
1.0	0.6346	0.0131	1.0	0.6714	0.01187

## D All Estimates

The following table reports the estimated bounds for all auctions  $t = 1, \dots, 39$ . The estimates are obtained using the standard bounds (Tight = no) both under risk neutrality ( $\vec{\rho} = (0, 0, 0)$ ) and under risk aversion ( $\vec{\rho} = \vec{\rho}^*$ ), as well as using the tighter bounds (Tight = yes) under risk aversion. Estimates are bagged from 200 bootstrap runs.

Auction	$AvP_l^{Pre}$			$AvP_u^{Pre}$			$AvP_l^{Post}$			$AvP_u^{Pre}$			$AvP_l^{Ratio}$			$AvP_u^{Ratio}$		
	$\vec{\rho}$	$\vec{\rho}$	$\vec{\rho}^*$	0	$\vec{\rho}^*$	$\vec{\rho}^*$	$\vec{\rho}$	$\vec{\rho}^*$	$\vec{\rho}^*$	0	$\vec{\rho}^*$	$\vec{\rho}^*$	$\vec{\rho}$	$\vec{\rho}^*$	$\vec{\rho}^*$	$\vec{\rho}$	$\vec{\rho}^*$	$\vec{\rho}^*$
Tight	no	no	yes	no	no	yes	no	no	yes	no	no	yes	no	no	yes	no	no	yes
1	8.172	7.867	7.923	12.13	12.03	9.142	0.3588	0.05413	0.1095	4.32	4.212	1.329	0.03457	0.006545	0.01604	0.4228	0.4189	0.1946
2	10.01	7.39	7.518	14.01	12.41	10.03	2.823	0.2001	0.329	6.823	5.221	2.842	0.1699	0.01767	0.04087	0.4931	0.4696	0.2922
3	6.479	6.346	6.399	10.73	10.69	7.548	0.1805	0.04697	0.09984	4.426	4.395	1.249	0.02849	0.008664	0.01695	0.4879	0.4869	0.2198
4	10.12	9.919	9.983	14.76	14.72	11.56	0.248	0.0511	0.1154	4.896	4.849	1.694	0.02344	0.005056	0.01423	0.3803	0.3791	0.1731
5	17.17	11.48	11.65	20.08	18.22	15.95	7.385	1.699	1.864	10.3	8.431	6.167	0.4213	0.07411	0.09253	0.5241	0.473	0.373
6	8.643	7.421	7.491	12.73	12.17	9.363	1.37	0.148	0.2177	5.461	4.9	2.09	0.1451	0.01482	0.02994	0.4313	0.4099	0.2384
7	9.988	9.691	9.764	14.15	14.09	11.06	0.3591	0.06177	0.135	4.518	4.456	1.427	0.04001	0.01039	0.02323	0.3742	0.3718	0.1686
8	16.44	12.42	12.57	20.34	19.97	18.28	4.847	0.8284	0.9719	8.739	8.374	6.689	0.3134	0.05891	0.07346	0.44	0.4285	0.3558
9	15.14	10.13	10.23	18.68	16.52	13.94	6.731	1.724	1.827	10.27	8.118	5.531	0.3756	0.1046	0.1293	0.5533	0.5122	0.3996
10	7.468	7.327	7.388	12.01	11.99	8.653	0.1843	0.04327	0.1046	4.731	4.709	1.37	0.02638	0.00655	0.01548	0.4411	0.4404	0.2
11	5.221	4.938	5.732	10.94	10.85	8.153	0.3333	0.04986	0.8445	6.049	5.966	3.265	0.09021	0.01536	0.06488	0.5713	0.5683	0.3602
12	7.507	7.382	7.434	13.04	13.02	8.894	0.1477	0.02347	0.07528	5.677	5.662	1.535	0.01538	0.003939	0.01318	0.4534	0.4527	0.2048
13	6.197	6.11	6.153	10.56	10.54	7.317	0.1167	0.0296	0.07329	4.476	4.456	1.237	0.02476	0.006056	0.01365	0.4807	0.4792	0.211
14	7.035	6.604	6.684	11.95	11.83	8.08	0.4734	0.0424	0.123	5.386	5.272	1.519	0.0356	0.006355	0.01546	0.523	0.521	0.2412
15	9.294	8.656	8.717	14.72	14.51	10.45	0.7006	0.06293	0.1238	6.125	5.917	1.86	0.06099	0.007748	0.02052	0.4596	0.4528	0.2258
16	16.4	10.55	10.67	19.66	16.07	13.72	7.032	1.188	1.305	10.29	6.702	4.355	0.3965	0.06713	0.08548	0.5343	0.486	0.352
17	10.95	9.37	9.471	14.51	13.81	11.32	1.718	0.1382	0.2383	5.273	4.575	2.092	0.09143	0.008704	0.02507	0.4295	0.4187	0.2217
18	11.65	8.636	8.776	15.41	14.23	11.67	3.323	0.3054	0.4456	7.084	5.903	3.344	0.2322	0.02078	0.04236	0.4745	0.4429	0.2864
19	7.647	6.793	6.863	14.0	13.73	9.121	0.9424	0.08901	0.1588	7.299	7.023	2.416	0.104	0.01361	0.03071	0.5223	0.5129	0.2785
20	7.295	6.606	6.67	12.28	11.95	8.195	0.7458	0.05735	0.1209	5.727	5.4	1.646	0.07643	0.009266	0.02227	0.5169	0.5088	0.2546
21	7.468	7.045	7.132	12.85	12.74	8.712	0.4703	0.04721	0.1339	5.849	5.737	1.714	0.0431	0.005956	0.01666	0.5223	0.5191	0.261
22	9.043	8.47	8.498	13.02	12.46	9.595	0.661	0.08765	0.1164	4.634	4.08	1.213	0.07118	0.005925	0.01859	0.4614	0.4511	0.2134
23	9.876	8.623	8.67	14.03	13.27	10.14	1.509	0.2552	0.3022	5.659	4.905	1.774	0.1103	0.01805	0.034	0.4493	0.437	0.2334
24	14.75	11.98	12.12	19.45	18.82	16.1	3.363	0.5934	0.7338	8.063	7.432	4.714	0.2064	0.02824	0.04738	0.4233	0.4069	0.286
25	17.76	12.6	12.67	20.42	19.76	17.56	8.654	3.499	3.57	11.32	10.66	8.453	0.4882	0.166	0.1819	0.5624	0.5392	0.4643
26	19.3	17.78	17.81	20.52	20.32	19.98	5.369	3.849	3.873	6.588	6.386	6.043	0.2624	0.1732	0.1786	0.3228	0.3201	0.3026
27	16.7	16.2	16.25	20.53	20.53	20.07	1.914	1.41	1.466	5.743	5.743	5.282	0.1421	0.1084	0.1169	0.2838	0.2838	0.2676
28	7.19	6.84	6.888	11.49	11.34	8.195	0.4004	0.0505	0.09793	4.695	4.552	1.405	0.04091	0.006956	0.01503	0.4738	0.4685	0.2207
29	5.718	5.121	5.219	9.345	9.035	7.096	0.6817	0.0854	0.1832	4.309	3.999	2.06	0.1264	0.01698	0.04533	0.4765	0.4647	0.2838
30	6.316	6.053	6.1	12.05	12.0	7.658	0.2959	0.03355	0.08076	6.031	5.981	1.638	0.03822	0.006484	0.01686	0.4901	0.4881	0.2311
31	9.951	9.852	9.871	20.53	20.53	14.38	0.11	0.01174	0.03014	10.69	10.69	4.543	0.03236	0.00388	0.005558	0.5183	0.5183	0.2847
32	18.23	16.59	16.65	20.52	20.49	20.11	6.697	5.057	5.112	8.982	8.953	8.579	0.3518	0.2798	0.2928	0.4508	0.449	0.4326
33	14.77	14.2	14.2	20.53	20.53	20.17	1.297	0.7214	0.7227	7.053	7.053	6.691	0.08257	0.05487	0.05726	0.3542	0.3542	0.3395
34	9.085	8.502	8.563	12.92	12.67	9.834	0.66	0.07719	0.1374	4.491	4.249	1.409	0.0723	0.00896	0.02675	0.4202	0.4112	0.2163
35	10.01	7.453	7.538	14.18	12.66	9.867	2.754	0.1969	0.2817	6.928	5.407	2.611	0.2379	0.02336	0.04362	0.5115	0.4758	0.3022
37	13.71	13.6	13.61	20.53	20.53	17.85	0.8014	0.6912	0.7018	7.625	7.625	4.948	0.09485	0.07648	0.08049	0.3719	0.3719	0.3399
37	14.55	10.92	10.97	18.13	15.82	13.52	4.615	0.9799	1.037	8.192	5.886	3.582	0.2914	0.06444	0.07839	0.4641	0.4128	0.2925
38	17.76	13.98	14.02	20.3	18.95	16.56	6.748	2.966	3.011	9.293	7.935	5.552	0.3847	0.1476	0.1674	0.4639	0.4435	0.3727
39	13.9	13.87	13.87	20.53	20.53	20.21	0.2226	0.1934	0.1957	6.857	6.857	6.536	0.01449	0.0127	0.01315	0.3344	0.3344	0.3102



The following table reports the standard errors of the estimated bounds for all auctions  $t = 1, \dots, 39$ . These are bootstrap standard errors obtained from 200 bootstrap estimates.

	$A_{v_i} P_i^{pre}$ (se)			$A_{v_i} P_i^{post}$ (se)			$A_{v_i} P_i^{pre}$ (se)			$A_{v_i} P_i^{ratio}$ (se)			$A_{v_i} P_i^{ratio}$ (se)			
	$\bar{p}$ Tight	$\bar{u}$ no	$\bar{p}^*$ no	$\bar{p}^*$ yes	$\bar{u}$ no	$\bar{p}^*$ no	$\bar{p}^*$ yes	0 no	$\bar{p}^*$ no	$\bar{p}^*$ yes	$\bar{u}$ no	$\bar{p}^*$ no	$\bar{p}^*$ yes	$\bar{u}$ no	$\bar{p}^*$ no	$\bar{p}^*$ yes
Auc.																
1	0.07566	0.01181	0.02179	0.08168	0.07566	0.01181	0.02179	0.09695	0.06773	0.08168	0.007253	0.001337	0.002659	0.004027	0.00316	0.007178
2	0.7328	0.1567	0.1601	0.4815	0.7328	0.1567	0.1601	0.7697	0.481	0.4815	0.03016	0.00746	0.009033	0.01841	0.01486	0.02212
3	0.04731	0.006673	0.02196	0.07348	0.04731	0.006673	0.02196	0.05003	0.03961	0.07348	0.007013	0.001237	0.002411	0.002128	0.001783	0.005443
4	0.1193	0.01534	0.03538	0.6489	0.1193	0.01534	0.03538	0.6489	0.6496	0.7077	0.005893	0.000838	0.004395	0.0004469	0.00431	0.01044
5	1.186	1.005	0.9861	1.156	1.186	1.005	0.9861	0.4007	0.9762	1.156	0.04541	0.04937	0.04717	0.00844	0.02271	0.04061
6	0.3225	0.1108	0.1114	0.2774	0.3225	0.1108	0.1114	0.3861	0.2774	0.2766	0.02574	0.005183	0.006249	0.01658	0.01328	0.0186
7	0.07125	0.004793	0.02845	0.0557	0.07125	0.004793	0.02845	0.07623	0.0557	0.1153	0.008323	0.002536	0.005638	0.005726	0.00522	0.01013
8	1.443	0.6371	0.6455	0.6937	1.443	0.6371	0.6455	0.2249	0.3219	0.6937	0.0602	0.03669	0.03546	0.004079	0.00949	0.03311
9	0.8858	0.8523	0.8477	0.8999	0.8858	0.8523	0.8477	0.7789	0.8497	0.8999	0.03274	0.04169	0.04144	0.01628	0.01965	0.03136
10	0.04771	0.005445	0.02105	0.0653	0.04771	0.005445	0.02105	0.04732	0.03879	0.0653	0.006575	0.0009595	0.002417	0.002318	0.002076	0.005967
11	0.07519	0.04182	0.7528	0.3547	0.07519	0.04182	0.7528	0.3658	0.3547	0.4316	0.01865	0.004673	0.01584	0.007493	0.006626	0.01736
12	0.03832	0.004492	0.1363	0.07963	0.03832	0.004492	0.1363	0.03796	0.03126	0.07963	0.004173	0.0007967	0.002089	0.002287	0.002082	0.006093
13	0.02173	0.00265	0.011	0.1345	0.02173	0.00265	0.011	0.09159	0.09035	0.1345	0.004371	0.0005851	0.001598	0.002681	0.002479	0.005291
14	0.09694	0.004358	0.0253	0.08515	0.09694	0.004358	0.0253	0.08938	0.06436	0.08515	0.006571	0.0005824	0.001731	0.002554	0.00203	0.004927
15	1.587	0.05224	0.02598	1.082	1.587	0.05224	0.02598	0.1597	0.09807	1.082	0.01355	0.0008741	0.002512	0.005585	0.003681	0.007718
16	1.397	0.7929	0.7787	0.9369	1.397	0.7929	0.7787	0.6998	0.7866	0.9369	0.05885	0.03542	0.03667	0.01127	0.01782	0.03857
17	0.5327	0.07982	0.09367	0.42	0.5327	0.07982	0.09367	0.5947	0.3923	0.42	0.02784	0.002482	0.004523	0.01407	0.01068	0.01938
18	0.8367	0.1924	0.1937	0.6158	0.8367	0.1924	0.1937	0.8573	0.6101	0.6158	0.04173	0.008407	0.01017	0.02246	0.01694	0.02545
19	0.2386	0.04612	0.04969	0.1694	0.2386	0.04612	0.04969	0.2593	0.1691	0.1694	0.02092	0.003449	0.00561	0.009209	0.006973	0.01417
20	0.169	0.01958	0.02897	0.1094	0.169	0.01958	0.02897	0.1947	0.1066	0.1094	0.01304	0.002162	0.003128	0.006065	0.004512	0.008877
21	0.08604	0.002189	0.02397	0.104	0.08604	0.002189	0.02397	0.1165	0.09649	0.104	0.007675	0.0004007	0.001941	0.003635	0.003049	0.006087
22	0.2328	0.05267	0.05444	0.2221	0.2328	0.05267	0.05444	0.4995	0.2221	0.2275	0.02177	0.00243	0.004187	0.01202	0.007611	0.01362
23	0.5366	0.1758	0.178	0.2222	0.5366	0.1758	0.178	0.5385	0.2099	0.2222	0.02894	0.006287	0.009081	0.01311	0.009263	0.01699
24	1.009	0.3266	0.3245	0.8778	1.009	0.3266	0.3245	0.5321	0.6639	0.8778	0.05279	0.01461	0.01439	0.01447	0.01829	0.03844
25	1.076	1.611	1.609	0.9424	1.076	1.611	1.609	0.2208	0.7734	0.9424	0.0252	0.0778	0.07505	0.005224	0.01682	0.03096
26	0.5763	1.17	1.167	0.5357	0.5763	1.17	1.167	0.02856	0.3083	0.5357	0.02342	0.05265	0.05155	0.0005476	0.005243	0.01249
27	0.6897	0.8188	0.8134	0.286	0.6897	0.8188	0.8134	0.007841	0.007841	0.286	0.03692	0.04881	0.04838	6.378e - 5	6.378e - 5	0.0005623
28	0.08098	0.009253	0.01471	0.2548	0.08098	0.009253	0.01471	0.2548	0.2404	0.2608	0.007669	0.0009001	0.001406	0.004086	0.002778	0.005941
29	0.1612	0.02424	0.1375	0.588	0.1612	0.02424	0.1375	0.588	0.5813	0.6446	0.02351	0.00307	0.006667	0.01479	0.01217	0.02064
30	0.05556	0.003361	0.01026	0.06843	0.05556	0.003361	0.01026	0.06843	0.05956	0.1428	0.006961	0.0005597	0.001983	0.002858	0.002436	0.006444
31	0.06764	0.008575	0.02546	3.562e - 15	0.06764	0.008575	0.02546	3.562e - 15	3.562e - 15	0.04051	0.01541	0.002749	0.003742	2.226e - 16	2.226e - 16	0.004245
32	1.126	1.609	1.607	0.0497	1.126	1.609	1.607	0.0497	0.08568	0.2216	0.0478	0.06045	0.05819	0.001329	0.002878	0.00967
33	0.5466	0.4588	0.4597	3.562e - 15	0.5466	0.4588	0.4597	3.562e - 15	3.562e - 15	0.5466	0.02866	0.02911	0.02911	5.565e - 17	5.565e - 17	0.0067
34	0.2042	0.03559	0.04171	0.2227	0.2042	0.03559	0.04171	0.2227	0.1091	0.1167	0.01807	0.003063	0.006842	0.01096	0.00811	0.01243
35	0.7782	0.1538	0.1519	0.4147	0.7782	0.1538	0.1519	0.911	0.4147	0.4075	0.04372	0.01117	0.01158	0.0232	0.01535	0.0242
36	0.3352	0.3427	0.3447	3.562e - 15	0.3352	0.3427	0.3447	3.562e - 15	3.562e - 15	0.0	0.04112	0.03936	0.04129	1.113e - 16	1.113e - 16	5.565e - 17
37	1.155	0.5727	0.5706	0.8196	1.155	0.5727	0.5706	1.086	0.8021	0.8196	0.04612	0.03317	0.03048	0.02306	0.02257	0.03462
38	1.015	1.274	1.264	0.8786	1.015	1.274	1.264	0.4572	0.5848	0.8786	0.03635	0.07558	0.07305	0.006845	0.01433	0.03972
39	0.1306	0.1165	0.1184	0.0	0.1306	0.1165	0.1184	8.904e - 16	8.904e - 16	0.0	0.0089	0.007523	0.007841	2.783e - 16	2.783e - 16	5.565e - 17

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